

# An algorithmic toolkit for continuous set covering on networks

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# Set covering problem

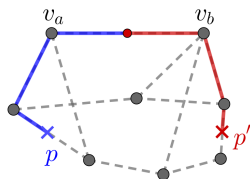
A classical NP-complete problem: Given a universal  $U$  and a family  $S$  of subsets of  $U$ , the problem asks the minimal number of subsets covering  $U$ .

$$\min_{s \in S} x_s \tag{1}$$

$$\sum_{e \in S} x_s \geq 1, \quad e \in U \tag{2}$$

$$x_s \in \{0, 1\}, \quad e \in S. \tag{3}$$

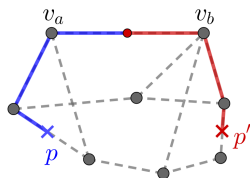
# Set covering problems on networks



(a) two facilities points

- ▶ Consider a network  $N = (V, A)$ ,
- ▶ Assume every edge is continuous, and its continuum is the union of points.
- ▶ The continuum of  $N$  is  $C(N)$ .

# Set covering problems on networks



(a) two facilities points  $p$  and  $p'$

- ▶ Let  $d(p_1, p_2)$  measure the shortest path distance between two points  $p_1$  and  $p_2$  in  $C(N)$ .
- ▶ Each point  $p \in C(N)$  can cover the points in  $C(N)$  with distance at most  $\delta$ .

# Extensions of set covering problems on networks

Here we use  $S$  to denote the set of facility locations, and  $U$  to denote the set of demands.

- ▶ *Discrete*: when  $U = V$  and  $S = V$ , reduced to the classical set covering problem (ILP or approximation algorithm).
- ▶ *Semi-continuous*: when either  $U = C(N)$  or  $S = C(N)$ , the problem is reducible to the classical set covering problem (tractable).
- ▶ *Continuous*: When both  $U = C(N)$  and  $S = C(N)$ , the *continuous set covering on networks*.

Some applications:

- ▶ locations of ambulance bases.
- ▶ surveillance cameras.
- ▶ routing servers in a network of computers.
- ▶ cranes for construction.
- ▶ aerial military medical evacuation facilities.
- ▶ aircraft alert sites for homeland defense.
- ▶ eVTOL safety landing sites.

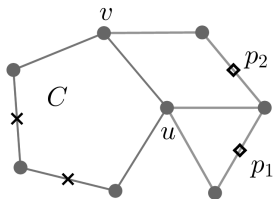
## Existing exact approach: discretization

- ▶ *discretization methods*: preprocessing procedures to reduce the problem to a tractable set covering problem.
- ▶ *finite dominating sets (FDS)*: finite subsets of candidate locations guaranteed to contain an optimal solution.



# Existing exact approach: discretization

An example: all edges have unit length and  $\delta = 2$ .



**Figure:** Two cycle coverage points with respect to a cycle  $C$  of five nodes

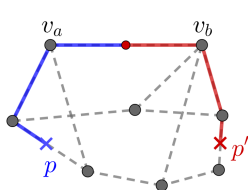
FDS: nodes and mid-points of edges.

## Existing exact approach: discretization

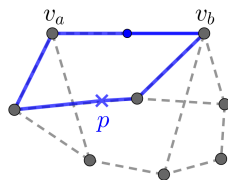
- ▶ *reduced problem*: semi-continuous, then further reduced to a discrete version.
- ▶ However, discretization methods rely on assumptions, e.g., edge lengths are natural numbers.

# Existing exact approach: MILP

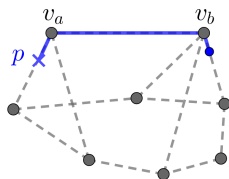
Covering conditions, an example:



(a) By two points



(b) By one point, through both ends



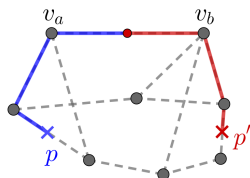
(c) By one point, through one end

# Existing exact approach: MILP

- ▶ The only existing MILP formulation is by Fröhlich et al., “Covering edges in network”.
- ▶ Basic assumption: edge length is at most  $\delta$ .
- ▶ MIP solvers cannot solve this MILP for moderate networks.

# Existing exact approach: MILP

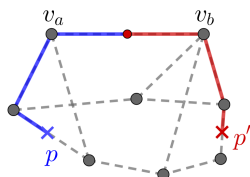
Basic ideas.



(a) By two points

- ▶ each edge can host a facility.
- ▶ an edge is covered, if the sum of available “cover range” from the left-end and the right-end is greater than the edge length.
- ▶ modeling the cover range between each pair of edges.

# Existing exact approach: MILP

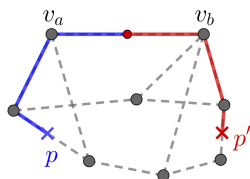


(a) By two points

Some comments:

- ▶ Each pair of edges are modeled, as in a complete graph. Networks are usually sparse!
- ▶ Some edges or nodes cannot contribute to cover, if the distance is large.
- ▶ Symmetry: if a facility is at a node, which incident edge host this node?

# New model of the covering condition



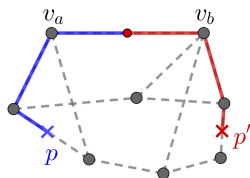
(a) By two points

- ▶ two kinds of facilities: facilities at nodes and facilities in edges.
- ▶ a point is covered by a covering path from a facility, the path length is at most  $\delta$ .

- ▶ Various preprocessing techniques: delimitation and modelling long edges.
- ▶ Two main new MILP models and some strengthening technique.
- ▶ an open-source implementation.



# Modeling the covering condition



(a) By two points

- ▶ the residual cover: given a node, the truncated length of covering paths.
- ▶ an edge is covered if the sum of residual cover from the left end and the right end is greater than the edge length.

# MILP Model 1

$$\min \sum_{f \in \mathcal{F}} y_f \quad (6a)$$

$$\text{s.t. } w_e \geq y_f \quad e \in E, f \in \mathcal{F}_c(e) \quad (6b)$$

$$w_e \leq \sum_{f \in \mathcal{F}_c(e)} y_f \quad e \in E \quad (6c)$$

$$x_v \geq 1 - \sum_{e \in E(v)} (1 - w_e) \quad v \in V \quad (6d)$$

$$x_v \leq w_e \quad v \in V, e \in E(v) \quad (6e)$$

$$y_{v_i'} + y_{e'} \leq 1 \quad e' \in E, i' \in \{a, b\} \quad (6f)$$

$$q_{e'} \leq l_{e'} y_{e'} \quad e' \in E \quad (6g)$$

$$l_e (1 - w_e) \leq r_{v_a} + r_{v_b} \quad e \in E \quad (6h)$$

$$x_v + \sum_{v' \in \mathcal{V}_p(v)} z_{vv'} + \sum_{(e', i') \in \mathcal{E}\mathcal{I}_p(v)} z_{ve'i'} = 1 \quad v \in V \quad (6i)$$

$$z_{vv'} \leq y_{v'} \quad v \in V, v' \in \mathcal{V}_p(v) \quad (6j)$$

$$z_{ve'i'} \leq y_{e'} \quad v \in V, (e', i') \in \mathcal{E}\mathcal{I}_p(v) \quad (6k)$$

$$r_v \leq M_v (1 - x_v) \quad v \in V \quad (6l)$$

$$r_v \leq M_{vv'} (1 - z_{vv'}) + \delta - d(v, v') \quad v \in V, v' \in \mathcal{V}_p(v) \quad (6m)$$

$$r_v \leq M_{ve'i'} (1 - z_{ve'i'}) + \delta - \tau_{ve'i'}(q_{e'}) \quad v \in V, (e', i') \in \mathcal{E}\mathcal{I}_p(v) \quad (6n)$$

$$y_f, w_e \in \{0, 1\} \quad f \in \mathcal{F}, e \in E \quad (6o)$$

$$x_v, z_{vv'}, z_{ve'i'} \in \{0, 1\} \quad v \in V, v' \in \mathcal{V}_p(v), (e', i') \in \mathcal{E}\mathcal{I}_p(v) \quad (6p)$$

$$q_{e'}, r_v \geq 0 \quad e' \in E, v \in V. \quad (6q)$$

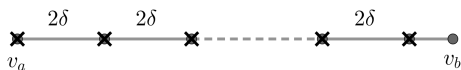
the reduction of the candidate space:

- ▶ **potential covers**: a set of edges and nodes in which a facility can possibly cover a given node.
- ▶ **complete covers**: a set of edges and nodes in which a facility can *always* cover a given edge.
- ▶ **partial covers**: further refinement of the potential covers.

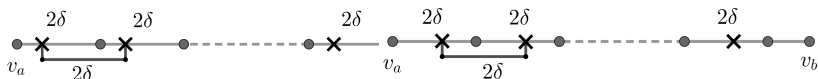
The previous modeling assumes short edges:  $l_e \leq \delta$ .

- ▶ First approach: subdivide long edges into small edges.
- ▶ Second approach: directly model the covering condition on long edge.

# Preprocessing: long edge modeling



(a) A facility is located at  $v_a$  ( $q_e = 0$ )



(b) A facility is located at the tail ( $0 < q_e \leq \hat{l}_e$ )

(c) No facility is located at the tail ( $\hat{l}_e < q_e \leq 2\delta$ )

Figure: Covering a long edge  $e = (v_a, v_b)$

Key observation: once the location of the left-most facility is determined, other facility locations are determined.

Modify MILP model 1 for covering on long edges. We add specific variables and constraints, and other parts remain the same.

- ▶ Implementation is based on JuMP and written in Julia.
- ▶ Input: a network and a cover radius  $\delta$ .
- ▶ Output: the number of facilities and locations.
- ▶ Algorithmic options:
  - ▶ EF: Covering edges in network.
  - ▶ F0: MILP model 1 without delimitation.
  - ▶ F: MILP model 1 with delimitation.
  - ▶ SF: MILP model 1 with delimitation and some valid inequalities.
  - ▶ RF: MILP model 2.
  - ▶ SFD: semi-continuous model, SF with facilities located at nodes.

- ▶ **Kgroup**: It consists of 23 prize-collecting Steiner tree problem instances, designed to have a local structure somewhat similar to street maps.
- ▶ **City**: It consists of real data of 9 street networks for some German cities.
- ▶ **Random**: It consists of 24 random networks instances generated by the package “Networkx”



- ▶ the relative dual gap is defined as:

$$\sigma := \frac{\bar{v} - \underline{v}}{\bar{v}},$$

where  $\bar{v}$  is an upper-bound and  $\underline{v}$  is a lower-bound.

- ▶ the relative primal bound

$$v_r := \frac{\bar{v}}{n_{sd}},$$

- ▶  $t$ : the total running time in CPU seconds.
- ▶ S/A/T: the number of solved instances/ the number of affected instances/ the number of total instances in the benchmark.

# Experimental results I

Benchmark	Radius	EF				FO			
		time	$\sigma(\%)$	$v_r(\%)$	S/A/T	time	$\sigma(\%)$	$v_r(\%)$	S/A/T
city	Small	1800.0	100.0%	100.0%	0/0/9	1801.7	56.8%	83.3%	0/3/9
	Large	1800.0	100.0%	100.0%	0/0/9	1800.9	42.3%	36.2%	0/6/9
Kgroup_A	Small	1800.0	100.0%	100.0%	0/0/11	1802.6	25.1%	85.0%	0/11/11
	Large	1800.0	100.0%	100.0%	0/0/11	139.2	14.7%	19.2%	7/11/11
Kgroup_B	Small	1800.0	100.0%	100.0%	0/0/12	1800.4	92.6%	98.8%	0/1/12
	Large	1800.0	100.0%	100.0%	0/0/12	1800.1	93.2%	86.6%	0/1/12
random_A	Small	1800.0	100.0%	100.0%	0/0/12	16.8	15.9%	54.8%	9/12/12
	Large	1800.0	100.0%	100.0%	0/0/12	0.2	25.5%	19.5%	12/12/12
random_B	Small	1800.0	100.0%	100.0%	0/0/12	1317.6	36.4%	63.3%	1/12/12
	Large	1800.0	100.0%	100.0%	0/0/12	154.4	26.0%	10.0%	11/12/12
all	Small	1800.0	100.0%	100.0%	0/0/56	625.8	37.4%	74.8%	10/39/56
	Large	1800.0	100.0%	100.0%	0/0/56	132.5	33.1%	25.9%	30/42/56

Table: Results for continuous models

# Experimental results II

Benchmark	Radius	F				SF			
		time	$\sigma(\%)$	$v_r(\%)$	S/A/T	time	$\sigma(\%)$	$v_r(\%)$	S/A/T
city	Small	1802.9	29.5%	62.2%	0/9/9	1801.3	30.1%	66.9%	0/9/9
	Large	1801.2	28.4%	21.7%	0/9/9	1800.9	29.1%	21.7%	0/9/9
Kgroup_A	Small	1803.0	33.1%	82.2%	0/11/11	1801.3	32.0%	80.6%	0/11/11
	Large	238.0	18.9%	19.1%	8/11/11	300.8	19.0%	19.1%	8/11/11
Kgroup_B	Small	1800.6	80.8%	240.5%	0/12/12	1801.4	79.7%	191.9%	0/12/12
	Large	1800.4	85.1%	80.5%	0/12/12	1800.7	85.9%	77.3%	0/12/12
random_A	Small	20.2	16.5%	54.3%	9/12/12	16.1	17.1%	54.9%	9/12/12
	Large	0.3	25.5%	19.5%	12/12/12	0.2	10.4%	17.9%	12/12/12
random_B	Small	1574.2	38.8%	64.9%	1/12/12	1501.2	40.0%	67.5%	1/12/12
	Large	220.5	19.9%	10.3%	9/12/12	175.7	18.8%	10.0%	11/12/12
all	Small	675.0	35.2%	86.2%	10/56/56	637.6	35.5%	83.6%	10/56/56
	Large	163.0	30.2%	23.6%	29/56/56	160.9	24.9%	22.8%	31/56/56

Table: Results for continuous models

# Experimental results III

Benchmark	Radius	RF			
		time	$\sigma$ (%)	$v_r$ (%)	S/A/T
city	Small	1804.4	16.2%	54.1%	0/9/9
	Large	1801.5	25.8%	21.3%	0/9/9
Kgroup_A	Small	1622.6	21.5%	77.5%	1/11/11
	Large	158.9	19.2%	19.3%	8/11/11
Kgroup_B	Small	1800.9	59.1%	154.2%	0/12/12
	Large	1800.6	75.5%	63.3%	0/12/12
random_A	Small	15.9	8.1%	54.3%	9/12/12
	Large	0.3	26.6%	19.8%	12/12/12
random_B	Small	1304.3	38.5%	63.8%	1/12/12
	Large	190.2	19.8%	11.2%	9/12/12
all	Small	604.9	23.7%	75.4%	11/56/56
	Large	146.6	29.2%	22.8%	29/56/56

Table: Results for continuous models

Code: <https://github.com/lidingxu/cflg>

Arxiv: Mercedes Pelegrín and Liding Xu, “Continuous Covering on Networks: Strong Mixed Integer Programming Formulations”

- ▶ New preprocessing and MILP models for continuous set-covering on networks.
- ▶ An open source implementation.